### Sparse Representations: from Source Separation to Compressed Sensing

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UNE UNITÉ DE RECHERCHE À LA POINTE DES SCIENCES ET DES TECHNOLOGIES DE L'INFORMATION ET DE LA COMMUNICATION

> Ecole d'été en Traitement du Signal Peyresq, Juillet 2009



#### Structure of the course

- Part I: Overview
- Part II: Algorithms, complexity & convergence
  - Lp minimization
  - Greedy Algorithms
- Part III: Recovery, stability, robustness
  - ♦ Null Space Properties and Lp minimization
  - Exact Recovery Condition and greedy algorithms
  - \* Restricted Isometry Constants, stability and robustness
- Part IV: Compressed Sensing and Random Matrices



#### Introduction

Compression
Adaptive
representation
Feature extraction
Kernel methods
(SVM ...)

Sparsity = old concept! (wavelets, ...)



Denoising
Blind source
separation
Compressed
sensing

Natural / traditional role:

Sparsity = low cost (bits, computations, ...)

direct goal

Novel indirect role

Sparsity = prior knowledge Tool for inverse problems



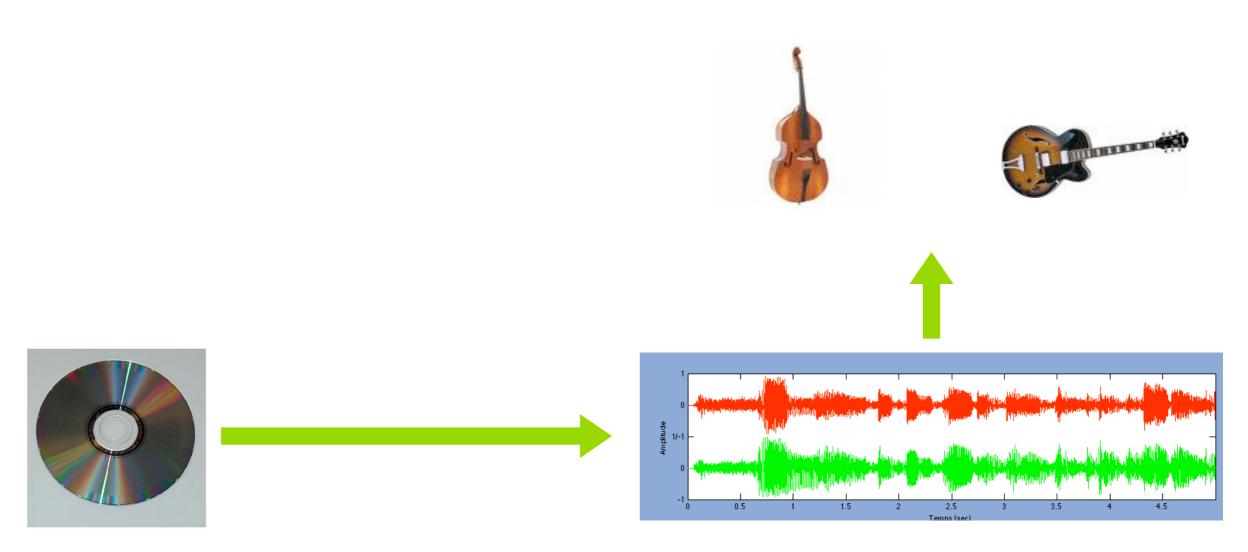
#### Overview

- Introduction : source separation and inverse problems
- Sparse decomposition algorithms
  - + LI minimisation
  - Matching Pursuits
- Provably good algorithms to recover sparse representations
- Compressed sensing & random sampling



## « Blind » Audio Source Separation

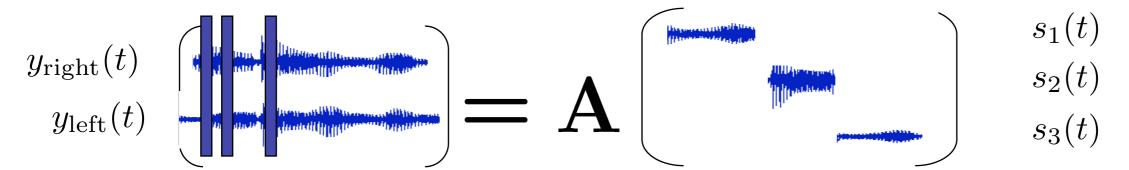
« Softly as in a morning sunrise »



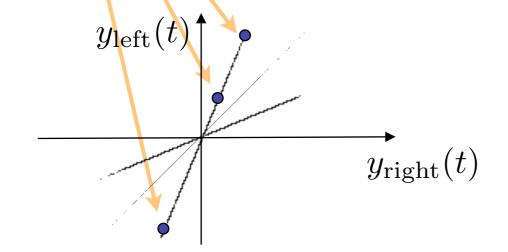


#### Blind Source Separation

Mixing model: linear instantaneous mixture



• Source model: if disjoint time-supports ...

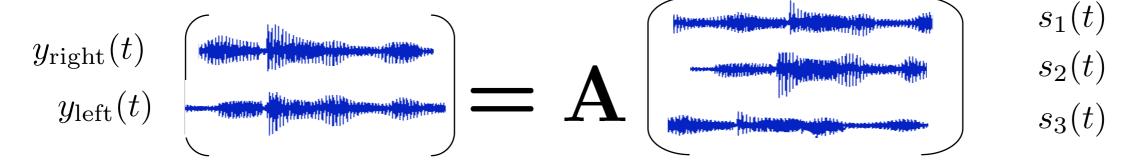


- ... then clustering to :
- 1- identify (columns of) the mixing matrix
- 2- recover sources

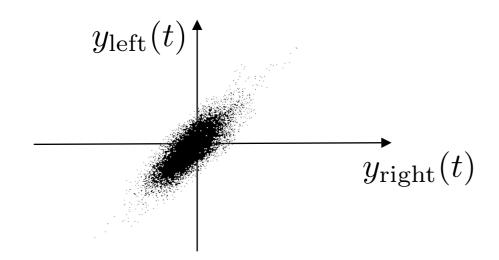


#### Blind Source Separation

Mixing model: linear instantaneous mixture



• In practice ...

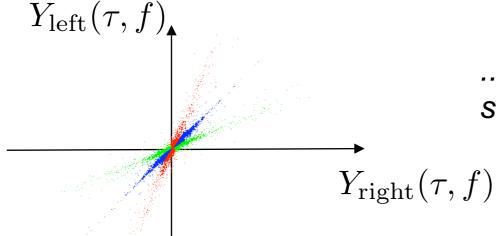


### Time-Frequency Masking

Mixing model in the time-frequency domain

$$egin{array}{ll} Y_{ ext{right}}( au,f) \ Y_{ ext{left}}( au,f) \end{array} = \mathbf{A} \ \mathbf{S}( au,f) \end{array}$$

And "miraculously" ...



... time-frequency representations of audio signals are (often) **almost disjoint**.

# Sparse decomposition algorithms

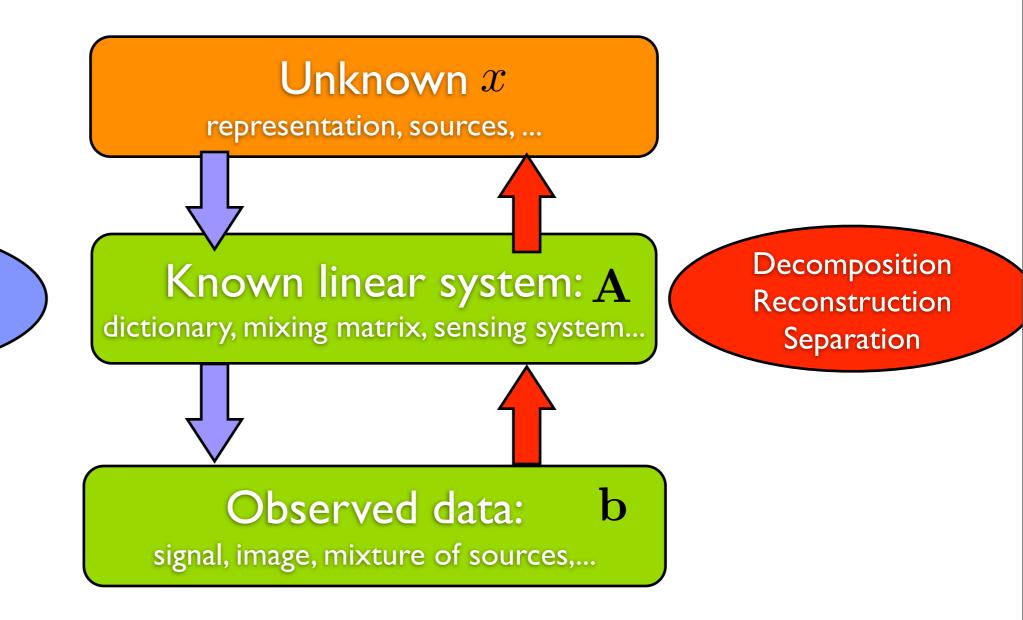


#### Vocabulary

Forward

linear model

 $\mathbf{b} \approx \mathbf{A} x$ 





### Sparsity and III-Posed Inverse Problems

• III-posedness if more unknowns than equations

$$\mathbf{A}x_0 = \mathbf{A}x_1 \not\Rightarrow x_0 = x_1$$

- Uniqueness of sparse solutions:
  - $\bullet$  if  $x_0, x_1$  are "sufficiently sparse",
  - + then  $\mathbf{A}x_0 = \mathbf{A}x_1 \Rightarrow x_0 = x_1$
- Recovery with practical algorithms
  - ◆ Thresholding, Matching Pursuits, Lp minimization p<=1,...</p>



#### Overall compromise

Approximation quality

$$\|\mathbf{A}x - \mathbf{b}\|_{2}$$

ullet Ideal sparsity measure :  $\ell^0$  "norm"

$$||x||_0 := \sharp \{n, \ x_n \neq 0\} = \sum_{n=1}^{\infty} |x_n|^0$$

Relaxed sparsity measure

$$||x||_p := (\sum_n |x_n|^p)^{1/p}$$



### Algorithms

Global optimization

Iterative greedy algorithms

Principle	$\min_{x} \frac{1}{2} \ \mathbf{A}x - \mathbf{b}\ _{2}^{2} + \lambda \ x\ _{p}^{p}$	
Tuning quality/sparsity	regularization parameter $\lambda$	
Variants	<ul> <li>choice of sparsity measure p</li> <li>optimization algorithm</li> <li>initialization</li> </ul>	

Principle

$$\min_{x} \frac{1}{2} \|\mathbf{A}x - \mathbf{b}\|_{2}^{2} + \lambda \|x\|_{p}^{p}$$

- **♦** Sparse representation  $\lambda \rightarrow 0$
- Sparse approximation  $\lambda > 0$



Principle

$$\min_{x} \frac{1}{2} \|\mathbf{A}x - \mathbf{b}\|_{2}^{2} + \lambda \|x\|_{p}^{p}$$

 $\lambda \to 0$ 

- ◆ Sparse representation
- Sparse approximation  $\lambda > 0$

Linear

$$p = 0 \qquad p = 2$$

Principle

$$\min_{x} \frac{1}{2} \|\mathbf{A}x - \mathbf{b}\|_{2}^{2} + \lambda \|x\|_{p}^{p}$$

 $\lambda \to 0$ 

- ◆ Sparse representation
- ◆ Sparse approximation

 $\lambda > 0$ 

NP-hard combinatorial

Linear

$$p = 0$$

$$p = 1$$

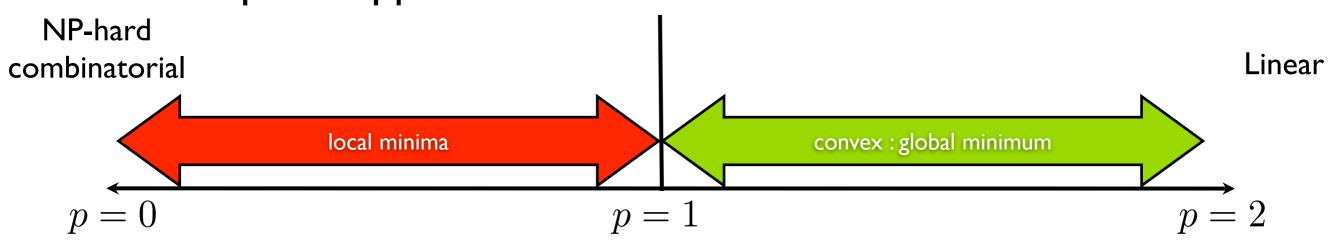
$$p=2$$

Principle

$$\min_{x} \frac{1}{2} \|\mathbf{A}x - \mathbf{b}\|_{2}^{2} + \lambda \|x\|_{p}^{p}$$

- Sparse representation
- $\lambda \to 0$
- ◆ Sparse approximation





Lasso [Tibshirani 1996], Basis Pursuit (Denoising) [Chen, Donoho & Saunders, 1999]

Linear/Quadratic programming (interior point, etc.)

Homotopy method [Osborne 2000] / Least Angle Regression [Efron &al 2002]

Iterative / proximal algorithms [Daubechies, Defrise, de Mol 2004, Combettes, & Pesquet 2008 ...]



Principle

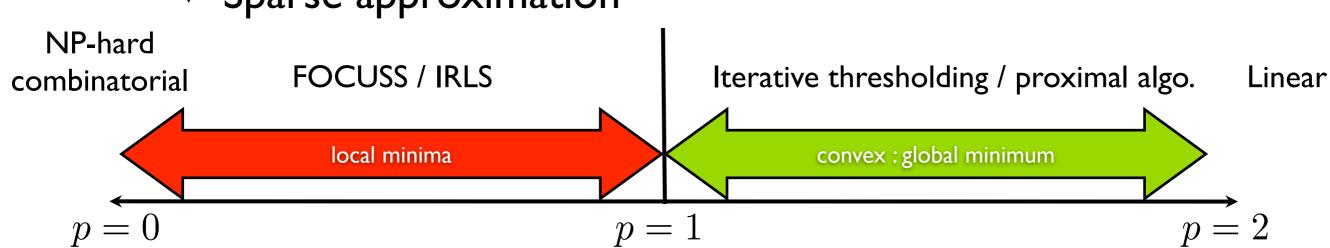
$$\min_{x} \frac{1}{2} \|\mathbf{A}x - \mathbf{b}\|_{2}^{2} + \lambda \|x\|_{p}^{p}$$

Sparse representation

 $\lambda \to 0$ 

◆ Sparse approximation

 $\lambda > 0$ 



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### Algorithms

Global optimization

Iterative greedy algorithms

Principle	$\min_{x} \frac{1}{2} \ \mathbf{A}x - \mathbf{b}\ _{2}^{2} + \lambda \ x\ _{p}^{p}$	iterative decomposition $\mathbf{r}_i = \mathbf{b} - \mathbf{A}x_i$ • select new components • update residual	
Tuning quality/sparsity	regularization parameter $\lambda$	stopping criterion (nb of iterations, error level,) $\ x_i\ _0 \geq k \qquad \ \mathbf{r}_i\  \leq \epsilon$	
Variants	<ul> <li>choice of sparsity measure p</li> <li>optimization algorithm</li> <li>initialization</li> </ul>	•selection criterion (stagewise) •update strategy (orthogonal)	



#### Main greedy algorithms

$$\mathbf{b} = \mathbf{A}x_i + \mathbf{r}_i$$

$$\mathbf{A} = [\mathbf{A}_1, \dots \mathbf{A}_N]$$

	Matching Pursuit	OMP	Stagewise
Selection	$\Gamma_i := \arg\max_n  \mathbf{A}_n^T \mathbf{r}_{i-1} $		$\Gamma_i := \{ n \mid  \mathbf{A}_n^T \mathbf{r}_{i-1}  > \theta_i \}$
	$\Lambda_i = \Lambda_{i-1} \cup \Gamma_i$	$\Lambda_i = \Lambda_{i-1} \cup \Gamma_i$	
Update	$x_i = x_{i-1} + \mathbf{A}_{\Gamma_i}^+ \mathbf{r}_{i-1}$	$x_i = \mathbf{A}_{\Lambda_i}^+ \mathbf{b}$	
	$\mathbf{r}_i = \mathbf{r}_{i-1} - \mathbf{A}_{\Gamma_i} \mathbf{A}_{\Gamma_i}^+ \mathbf{r}_{i-1}$	$\mathbf{r}_i =$	$\mathbf{b} - \mathbf{A}_{\Lambda_i} x_i$

MP & OMP: Mallat & Zhang 1993

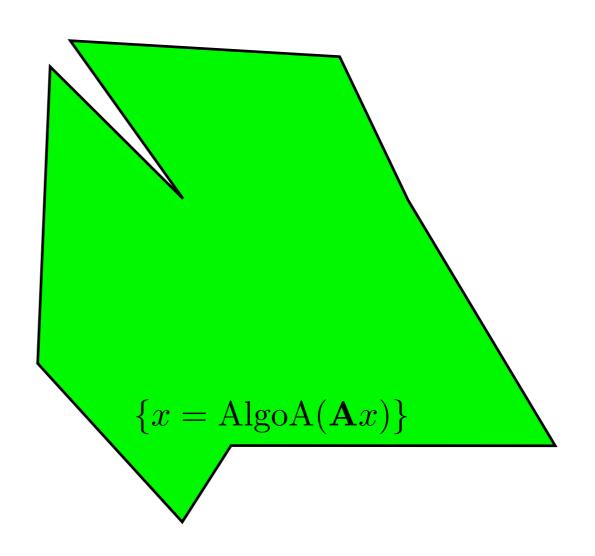
StOMP: Donoho & al 2006 (similar to MCA, Bobin & al 2006)



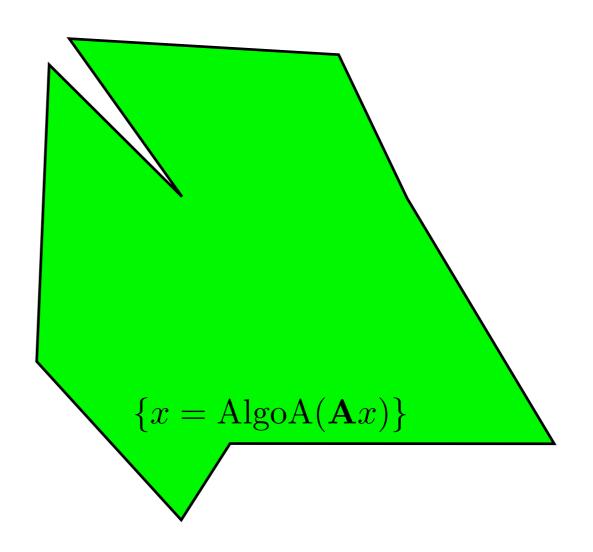
### Provably good algorithms



• Recoverable set for a given "inversion" algorithm

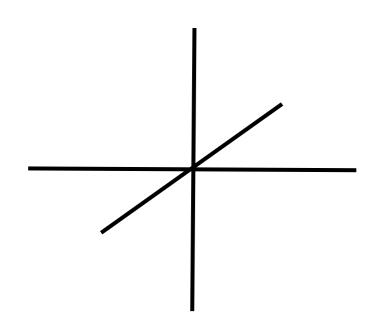


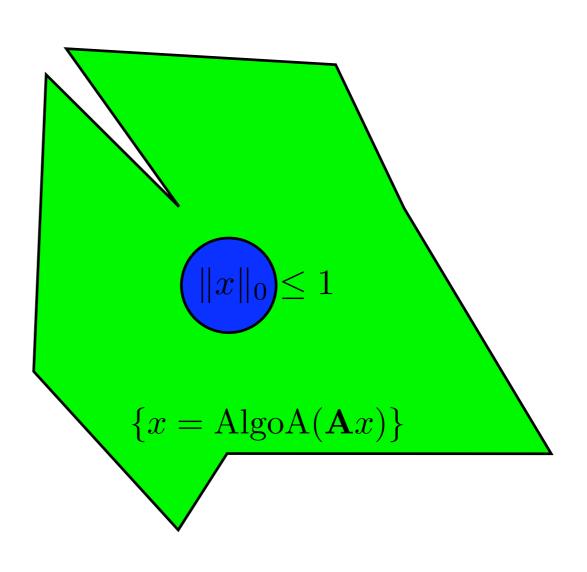
- Recoverable set for a given "inversion" algorithm
- Level sets of L0-norm





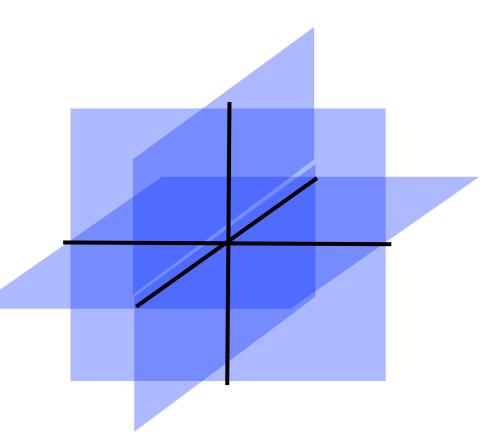
- Recoverable set for a given "inversion" algorithm
- Level sets of L0-norm
  - ◆ I-sparse

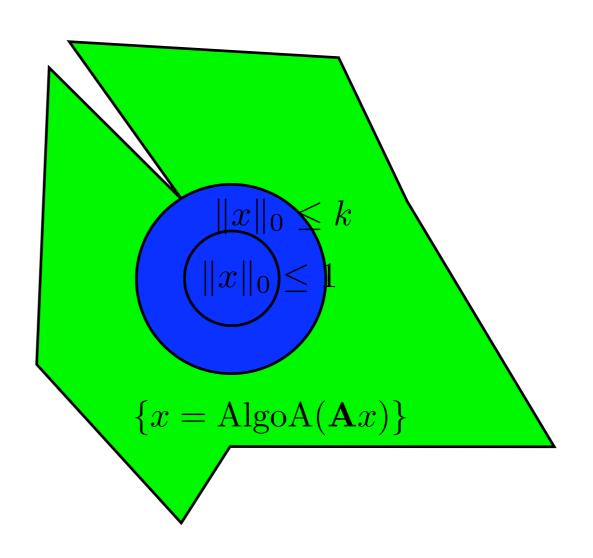






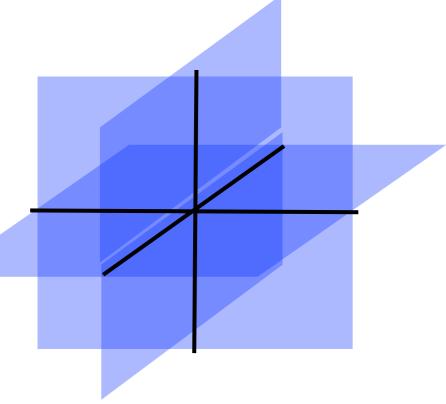
- Recoverable set for a given "inversion" algorithm
- Level sets of L0-norm
  - ◆ I-sparse
  - → 2-sparse

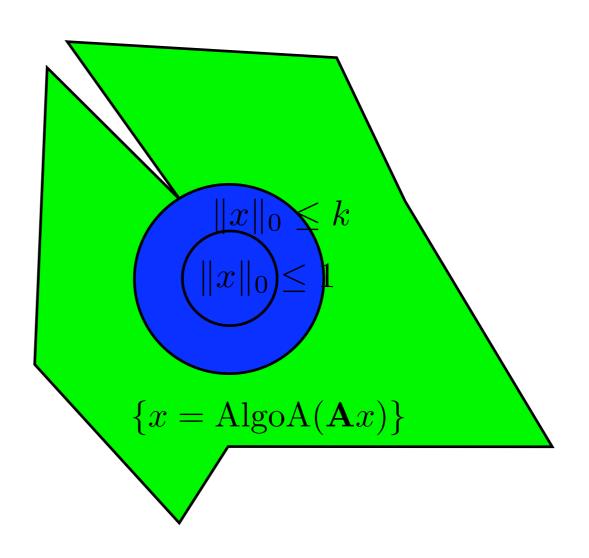






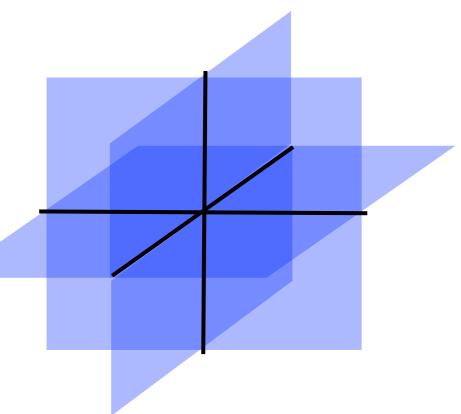
- Recoverable set for a given "inversion" algorithm
- Level sets of L0-norm
  - ◆ I-sparse
  - → 2-sparse
  - **→** 3-sparse ...

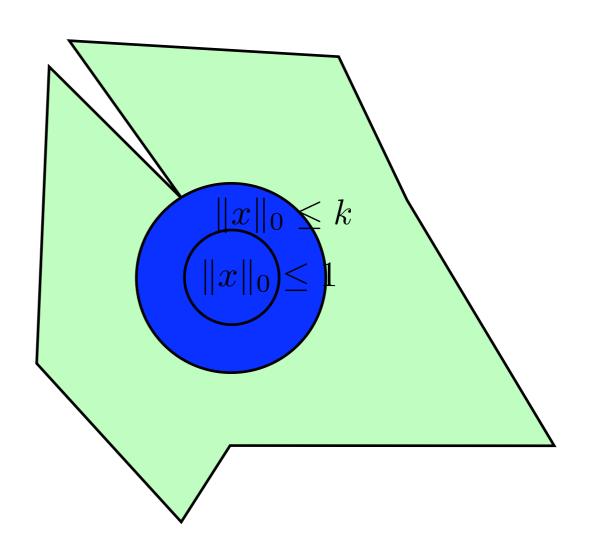






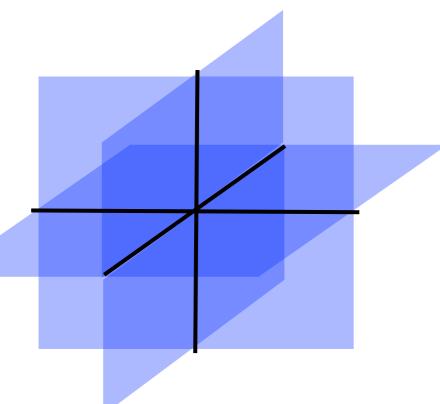
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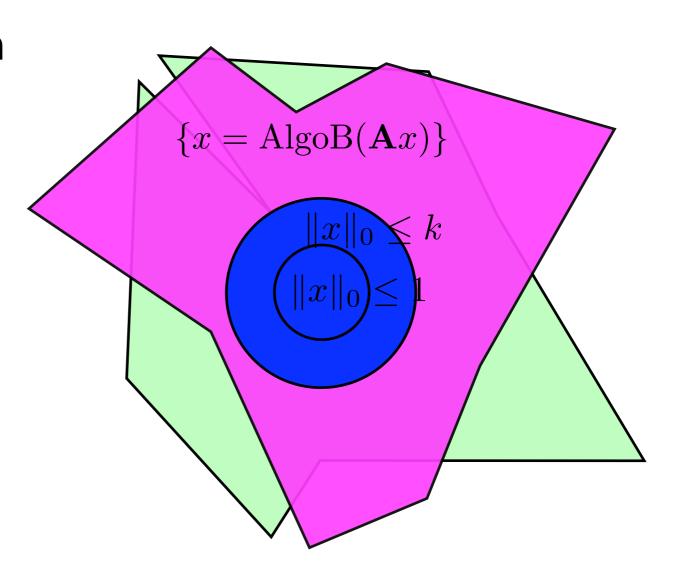






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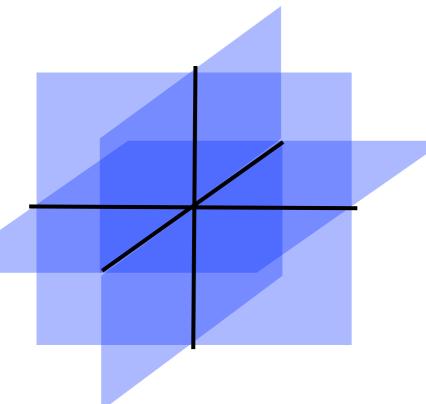


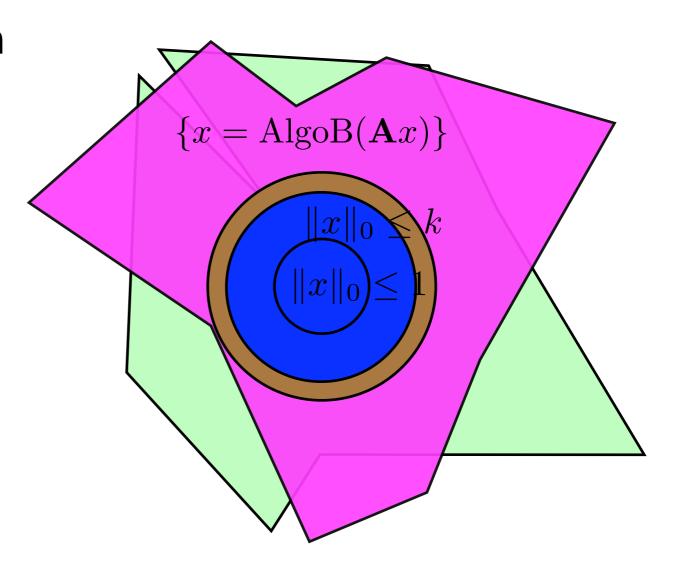




## Recovery analysis for inverse problem b = Ax

- Recoverable set for a given "inversion" algorithm
- Level sets of L0-norm
  - ◆ I-sparse
  - → 2-sparse
  - **→** 3-sparse ...

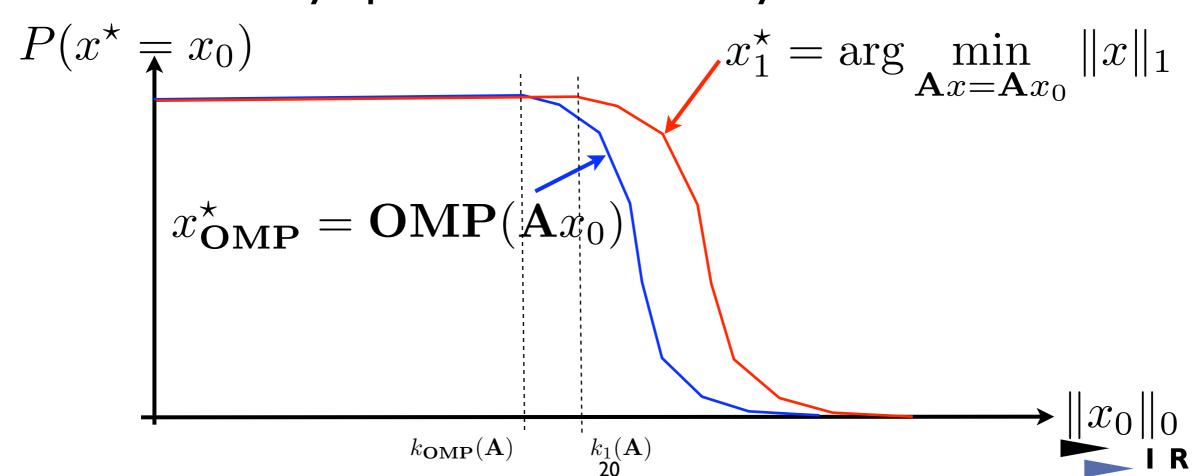






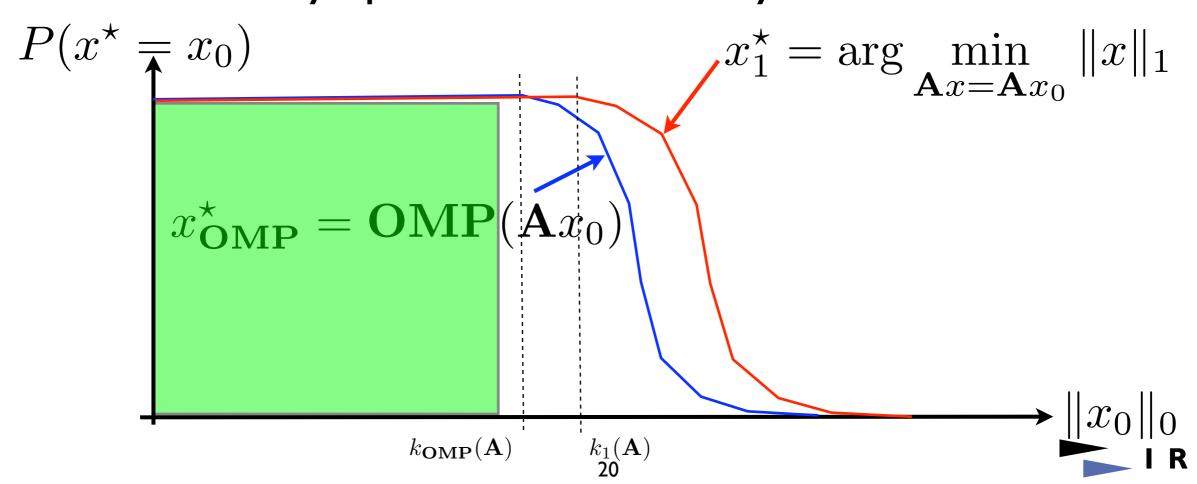
#### Empirical facts

- Highly sparse vectors: always recovered by
  - Orthonormal Matching Pursuit
  - ◆ Basis Pursuit (LI minimization)
- Relatively sparse vectors: likely to be recovered



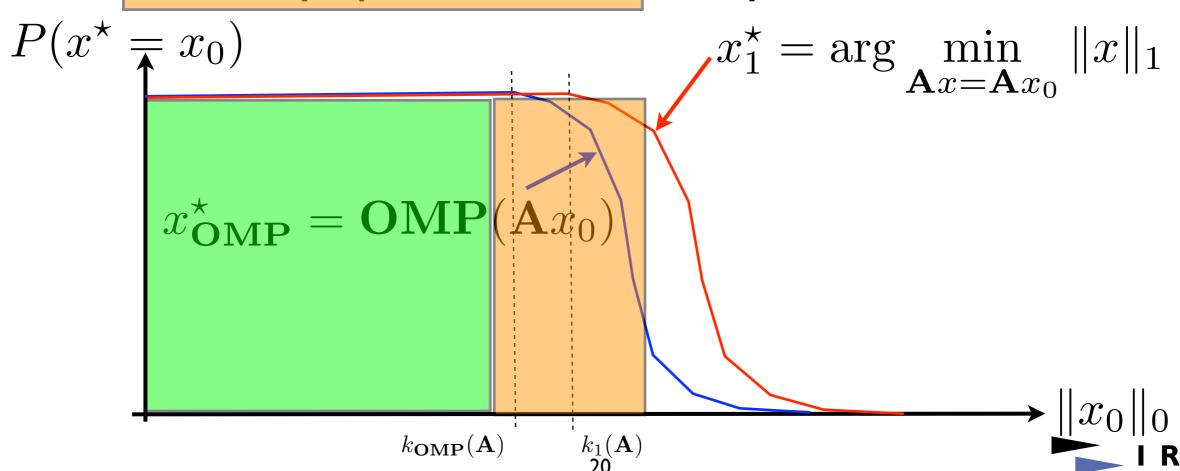
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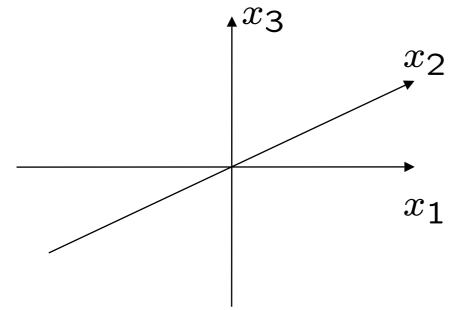


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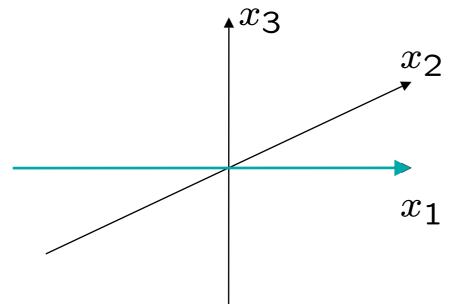


Brute force search



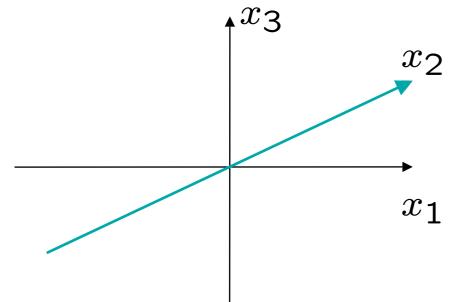
- Theorem (Davies et al, Natarajan)
  - It is NP-hard!
- Are there other more efficient alternatives?

Brute force search



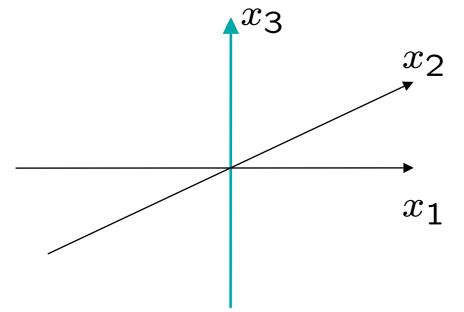
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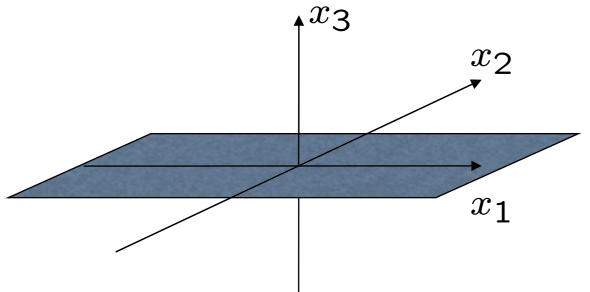
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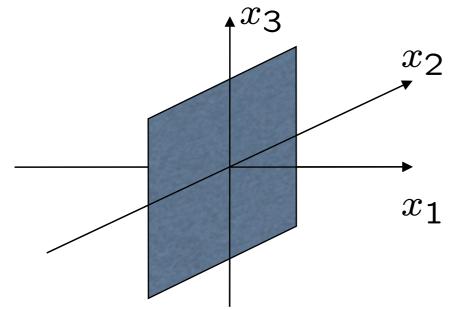
Brute force search



 $\min_{x} \|\mathbf{b} - \mathbf{A}x\|_{2}^{2} + \lambda \|x\|_{0}$ 

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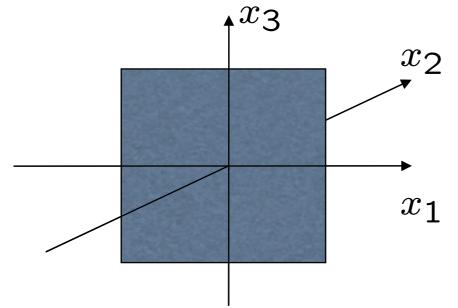
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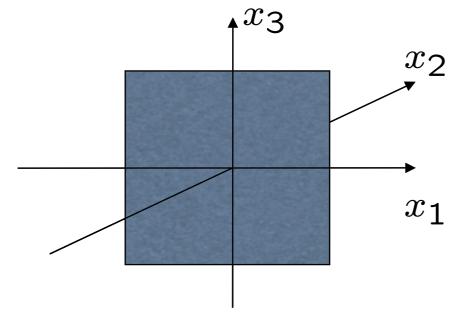
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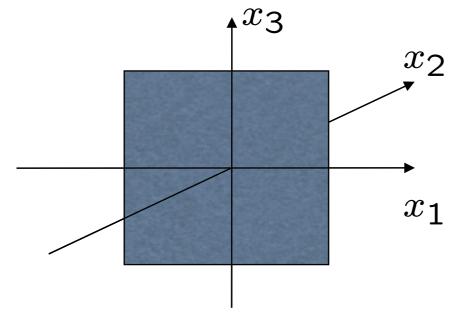
 $\min_{x} \|\mathbf{b} - \mathbf{A}x\|_{2}^{2} + \lambda \|x\|_{0}$ 

- Theorem (Davies et al, Natarajan)
  - It is NP-hard!
- Are there other more efficient alternatives?



$$\begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \end{bmatrix} \cdot \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \mathbf{b}$$

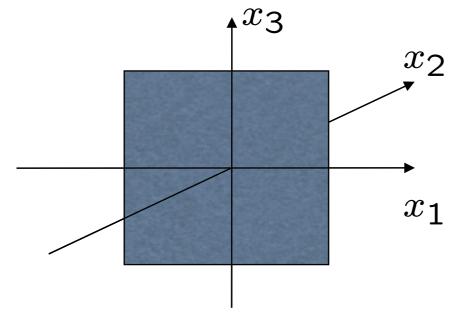
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$$\begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \end{bmatrix} \cdot \begin{pmatrix} x_1 \\ 0 \\ x_3 \end{pmatrix} = \mathbf{b}$$

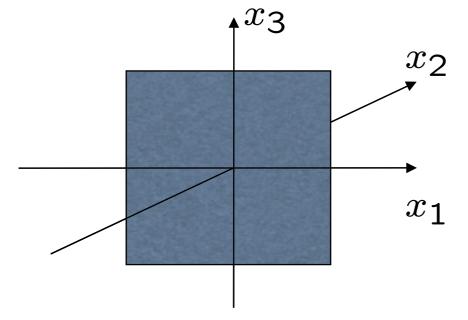
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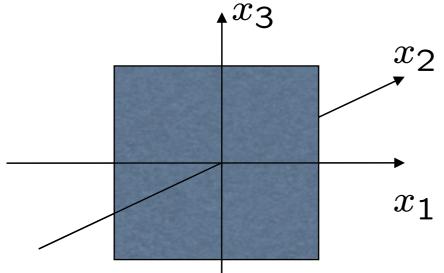
$$\begin{bmatrix} a_{11} & a_{13} \\ a_{21} & a_{23} \end{bmatrix} \cdot \begin{pmatrix} x_1 \\ x_3 \end{pmatrix} = \mathbf{b}$$

$$\begin{pmatrix} x_1 \\ x_3 \end{pmatrix} = \begin{bmatrix} a_{11} & a_{13} \\ a_{21} & a_{23} \end{bmatrix}^{-1} \cdot \mathbf{b}$$

- Theorem (Davies et al, Natarajan)
  - It is NP-hard!
- Are there other more efficient alternatives?



Brute force search



$$\left[\begin{array}{cc} a_{11} & a_{13} \\ a_{21} & a_{23} \end{array}\right] \cdot \left(\begin{array}{c} x_1 \\ x_3 \end{array}\right) = \mathbf{b}$$

$$\begin{pmatrix} x_1 \\ x_3 \end{pmatrix} = \begin{bmatrix} a_{11} & a_{13} \\ a_{21} & a_{23} \end{bmatrix}^{-1} \cdot \mathbf{b}$$

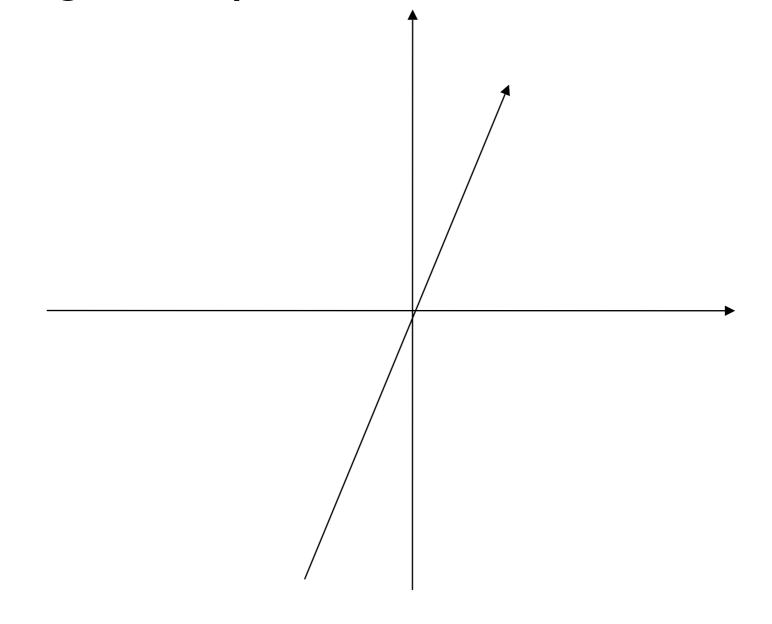
Many n-tuples to try!

- Theorem (Davies et al, Natarajan)
  - It is NP-hard!
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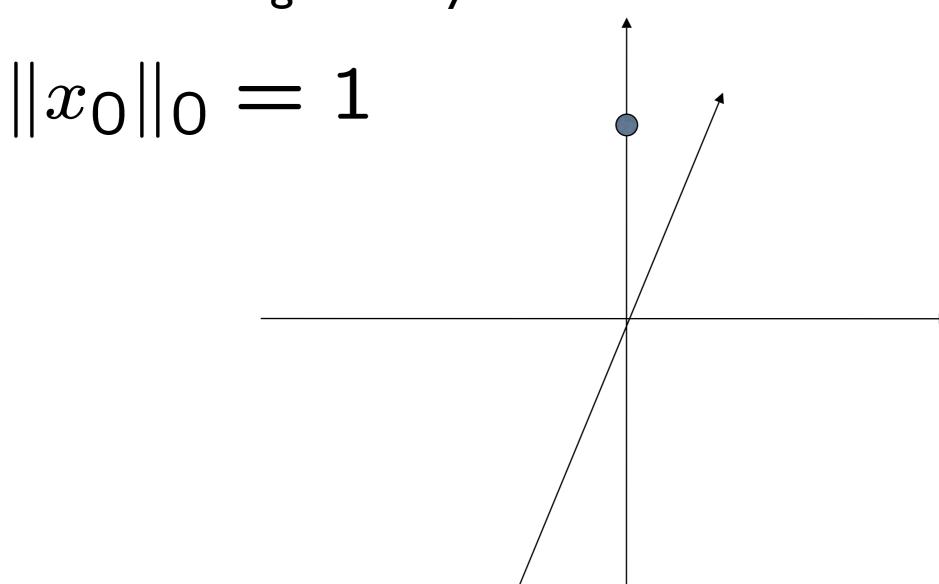
Some 3D geometry

$$\mathbf{b} = \mathbf{A}x_0$$



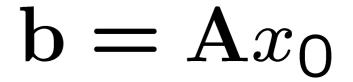
Some 3D geometry

$$\mathbf{b} = \mathbf{A}x_0$$



Some 3D geometry

$$||x_0||_0 = 1$$



$$\{\mathbf{A}x = \mathbf{A}x_0\}$$

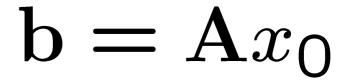


Some 3D geometry

$$||x_0||_0 = 1$$

Sparse solutions

$$\left\{ \|x\|_{1} \leq \|x_{0}\|_{1} \right\}$$



$$\left\{\mathbf{A}x = \mathbf{A}x_{\mathbf{0}}\right\}$$



Some 3D geometry

$$\mathbf{b} = \mathbf{A}x_0$$

$$||x_0||_0 = 1$$

Sparse solutions

$$\left\{ \|x\|_{1} \leq \|x_{0}\|_{1} \right\}$$

$$\{\mathbf{A}x = \mathbf{A}x_0\}$$



Some 3D geometry

$$\mathbf{b} = \mathbf{A}x_0$$

$$||x_0||_0 = 2$$

Sparse solutions

$$\left\{ \|x\|_{1} \leq \|x_{0}\|_{1} \right\}$$

$$\{\mathbf{A}x = \mathbf{A}x_0\}$$



Some 3D geometry

$$||x_0||_0 = 2$$

Sparse solutions

$$\left\{ \|x\|_{1} \leq \|x_{0}\|_{1} \right\}$$

 $\mathbf{b} = \mathbf{A}x_0$ 

$$x_0 = \arg\min_{x:\mathbf{A}x=\mathbf{b}} \|x\|_1$$

$$\{\mathbf{A}x = \mathbf{A}x_0\}$$

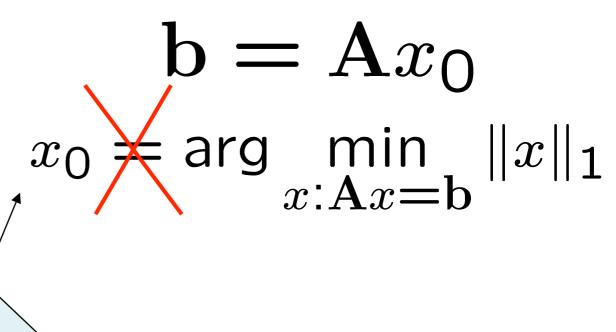


Some 3D geometry

$$||x_0||_0 = 3$$

Sparse solutions

$$\left\{ \|x\|_{1} \leq \|x_{0}\|_{1} \right\}$$



$$\{\mathbf{A}x = \mathbf{A}x_0\}$$



# Equivalence between L0, L1, OMP

• Theorem : assume that  $\mathbf{b} = \mathbf{A}x_0$ 

+ if 
$$||x_0||_0 \le k_0(\mathbf{A})$$
 then  $x_0 = x_0^*$ 

$$\star$$
 if  $||x_0||_0 \le k_1(\mathbf{A})$  then  $x_0 = x_1^{\star}$ 

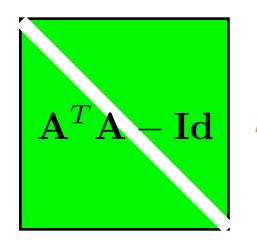
where 
$$x_p^{\star} = \arg\min_{\mathbf{A}x = \mathbf{A}x_0} \|x\|_p$$

- Donoho & Huo 01 : pair of bases, coherence
- Donoho & Elad, Gribonval & Nielsen 2003 : dictionary, coherence
- Tropp 2004: Orthonormal Matching Pursuit, cumulative coherence
- Candes, Romberg, Tao 2004: random dictionaries, restricted isometry constants



### State of the art tools to

estimate  $k_p(\mathbf{A})$ 



max over N(N-I) entries



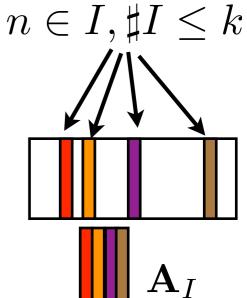




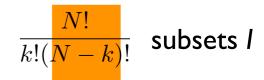
N unit columns

$$\|\mathbf{A}_n\|_2 = 1$$





max over



$$\delta_k := \sup_{\sharp I \le k, \ c \in \mathbb{R}^k} \left| \frac{\|\mathbf{A}_I c\|_2^2}{\|c\|_2^2} - 1 \right|$$

$$\widehat{k}(\mathbf{A}) = (1 + 1/\mu)/2$$

 $\mu = \mu(\mathbf{A}) := \max_{k \neq l} |\langle \mathbf{A}_k, \mathbf{A}_l \rangle|$ 

$$\delta_{2k_0} < 1$$

$$\delta_{2k_1} < \sqrt{2} - 1$$

#### (Cumulative) coherence

Low cost "Coarse / pessimistic" Common beliefs

#### **Isometry constants**

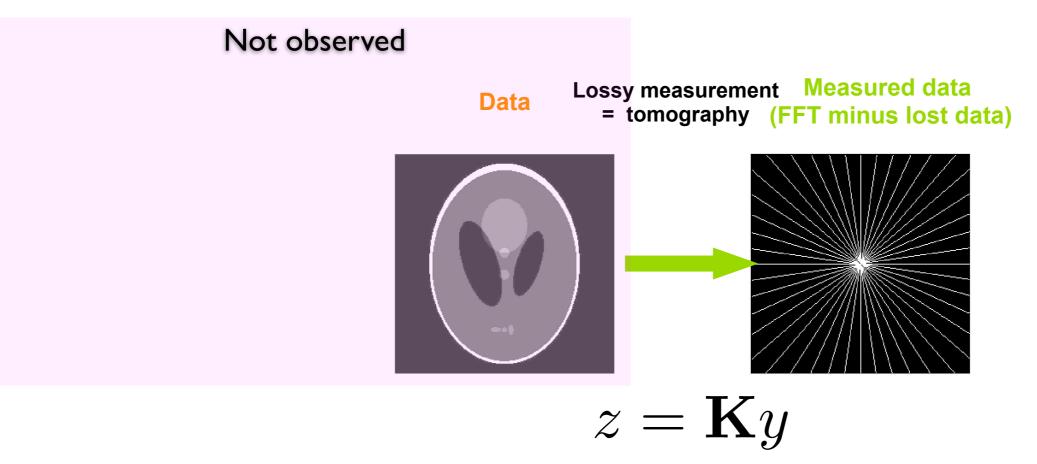
Hard to compute "Almost sharp"?



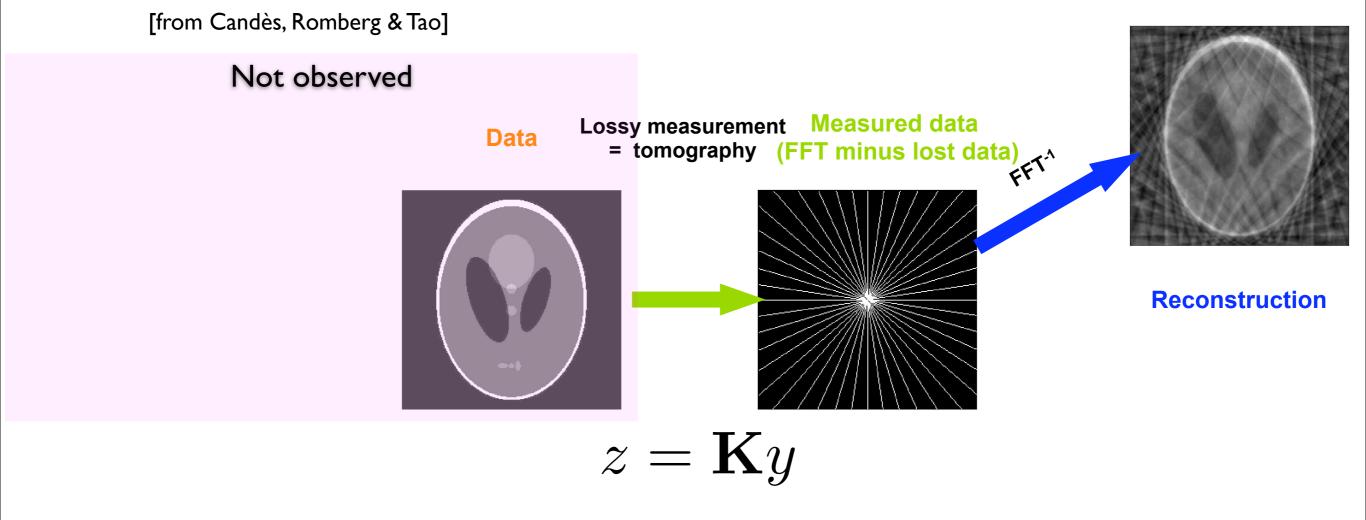


MRI from incomplete measures

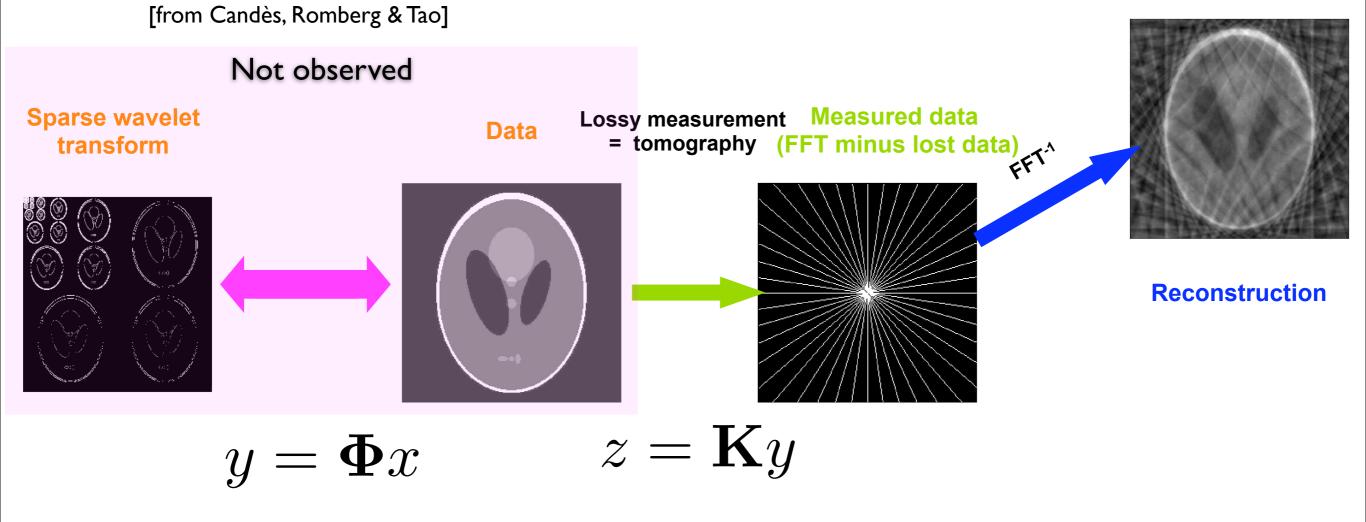
[from Candès, Romberg & Tao]



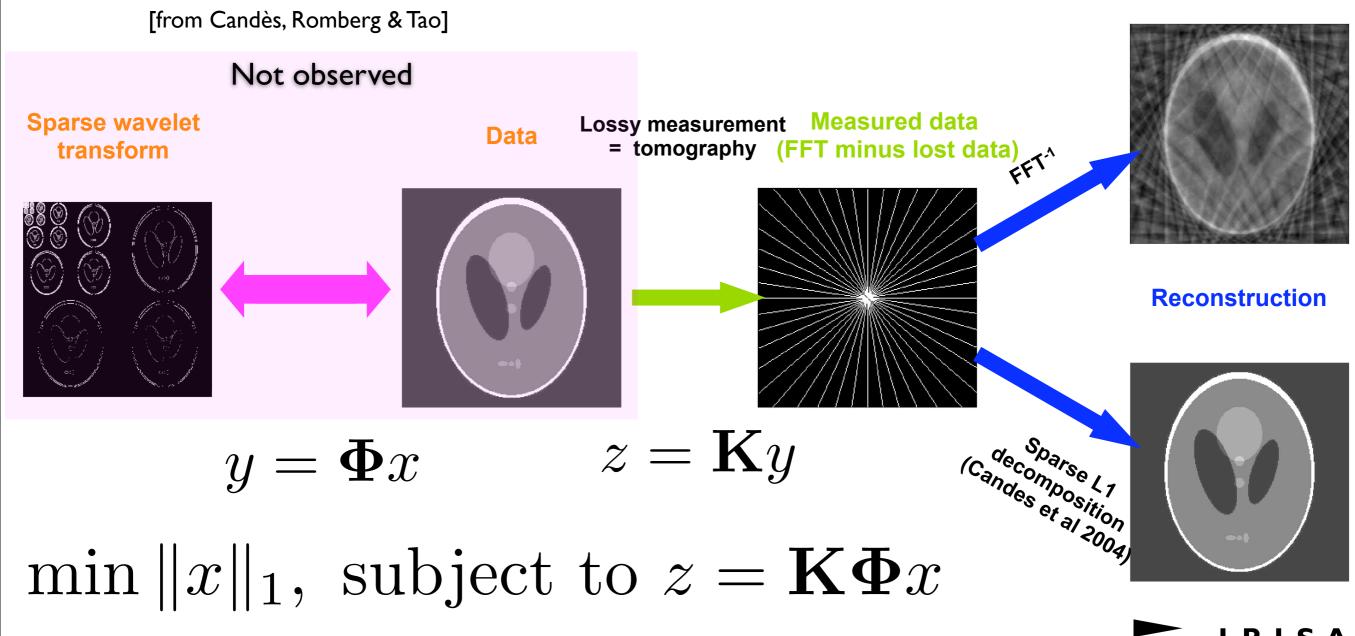
MRI from incomplete measures



MRI from incomplete measures

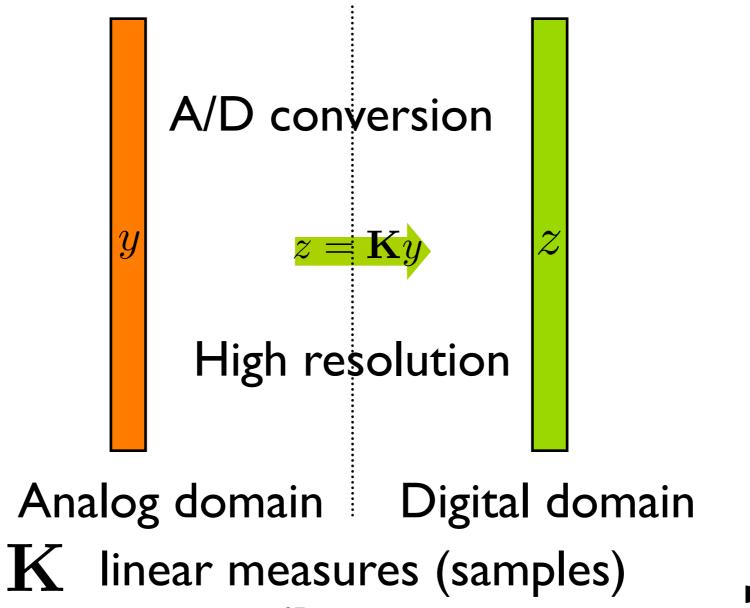


MRI from incomplete measures



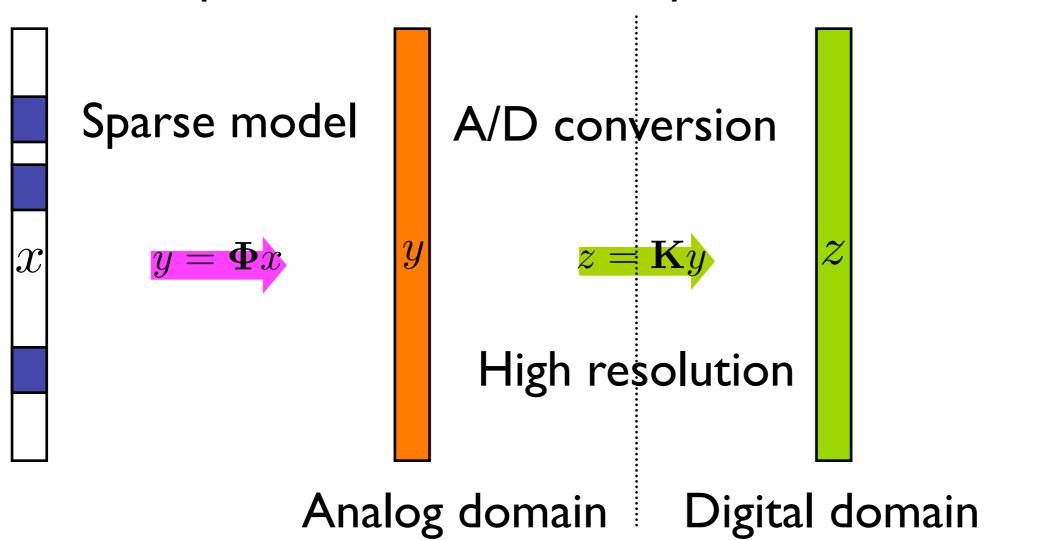
### Classical Shannon Sampling

« Sample first, think and compress afterwards »



### Classical Shannon Sampling

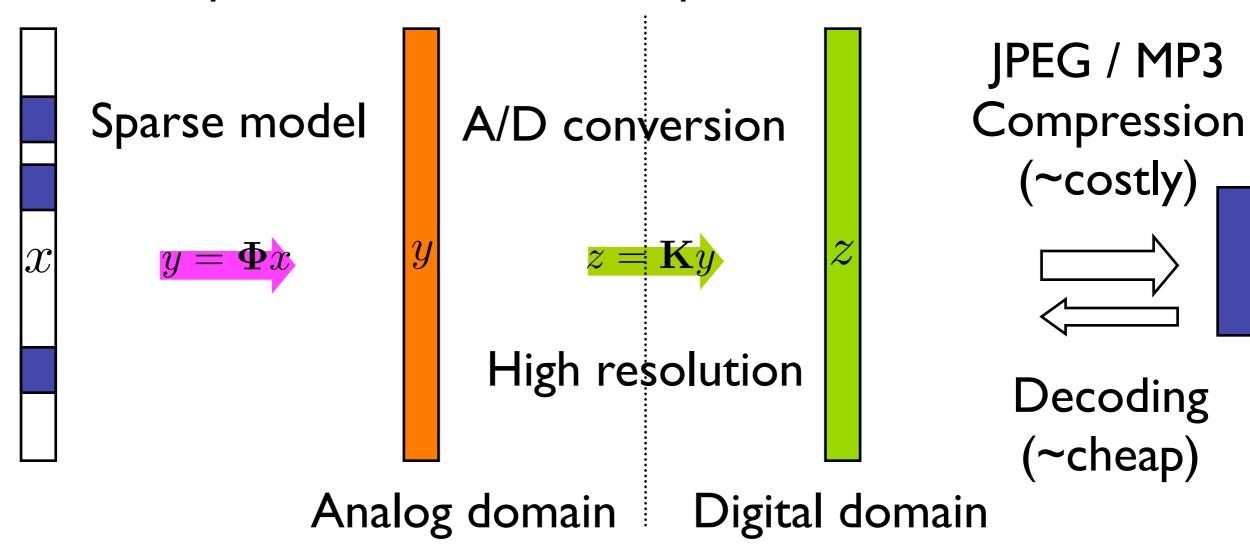
« Sample first, think and compress afterwards »





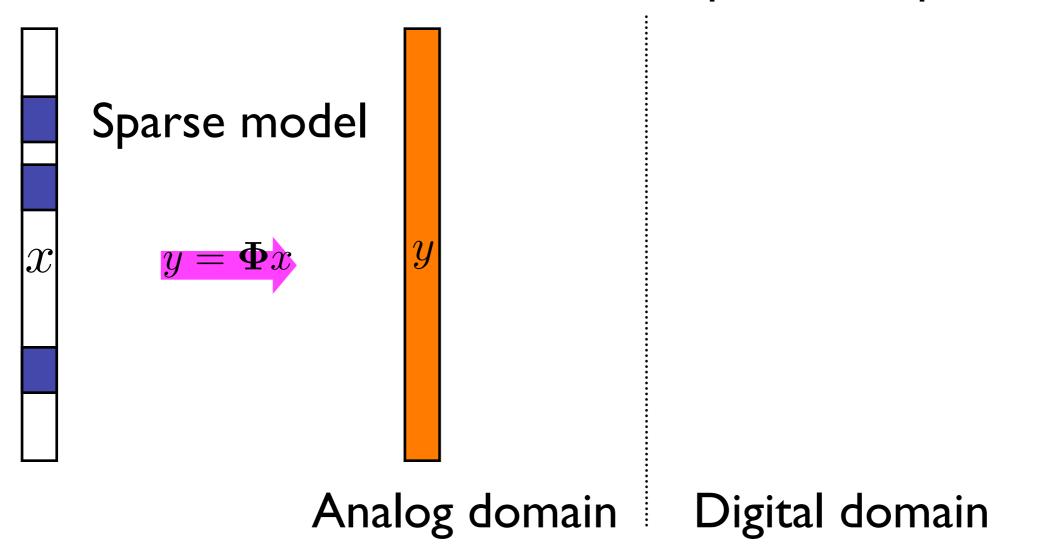
### Classical Shannon Sampling

« Sample first, think and compress afterwards »



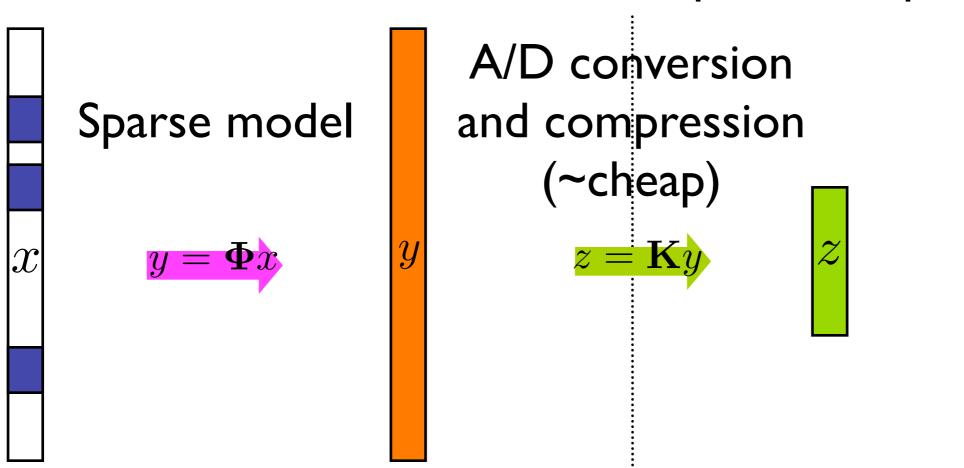


• First model the data, then sample & compress





• First model the data, then sample & compress



Analog domain

Digital domain



First model the data, then sample & compress

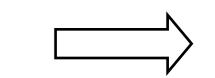
Sparse model

$$y = \Phi x$$

and compression (~cheap)

$$z = \mathbf{K}y$$

A/D conversion Sparse recovery (~costly)



Analog domain Digital domain

 $\min \|x\|_1$ , subject to  $z = \mathbf{K}\Phi x$ 



### Partial conclusions

- Sparsity helps solve ill-posed inverse problems (more unknowns than equations)
- Sparse approximation is NP-hard but efficient suboptimal algorithms (pursuits) exist
- If there is a sufficiently sparse solution, it is recovered by many of the heuristic algorithms
- This is the fundamental basis underlying the development of compressed sensing



### Structure of the course

- Part I: Overview
- Part II: Algorithms, complexity & convergence
  - Lp minimization
  - Greedy Algorithms
- Part III: Recovery, stability, robustness
  - ◆ Null Space Properties and Lp minimization
  - ◆ Exact Recovery Condition and greedy algorithms
  - \* Restricted Isometry Constants, stability and robustness
- Part IV: Dictionaries, Random Matrices and Compressed Sensing

